

# CALCULUS IN NON-LINEAR PROGRAMMING PROBLEM

By

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In the recent paper [1] Wilde has derived with the help of calculus, the necessary conditions for a relative minimum of non-linear differentiable objective function of non-negative values constrained by non-linear differentiable inequalities.

It is possible to establish the sufficiency of these conditions *by using only well known theorems of differential calculus* provided we assume further that the objective function is such that it can be transformed into a strictly convex function under given constraints. (Assuming sufficient differentiability, an anisotropically convex function may be transformed into a convex function by application of the theory of invariance of quadratic forms [4].) The results are expressed entirely in terms of partial derivatives.

According to Wilde a non-linear programming problem is to find a set of non-negative variables,  $x_j (j=1, 2, \dots, J)$  that satisfy a set of differentiable constraints:

$$f_i(x_j) \geq 0 \quad (i=1, 2, \dots, I) \quad (1)$$

and that also minimise a differentiable non-linear objective function  $g(x_i)$ .

Let  $f_i(x_j)$  represent  $I$  functions mapping  $J$  real variables  $x_j$  into real variables  $f_i$ . This transformation may be expressed by the following  $I$  equations in the  $I+J$  variables  $f_i$  and  $x_j$ .

$$f_i - f_i(x_j) = 0 \quad \left. \begin{matrix} i=1, \dots, I \\ j=1, \dots, J \end{matrix} \right\} \quad (2)$$

Let these equations have real solution  $(f_i^0, x_j^0)$  such that

$$f_i^0 - f_i^0(x_j^0) = 0 \quad (3)$$

The differentiability of  $f_i(x_j)$  guarantees that the  $I$  functions  $f_i = f_i(x_j)$  regarded as functions of all  $I+J$  independent variables are differentiable in the neighbourhood of the point  $(f_i^0, x_j^0)$ .

From among these  $I+J$  independent variables, let us select any  $I$  say  $x_l (l=1, \dots, p)$  and  $f_q (q=p+1, \dots, I)$  where  $1 \leq p \leq I$

If Jacobian

$$|J(f_i^0, x_j^0)| \neq 0 \quad (4)$$

then the set of equations (2) can be solved for  $I$  variables  $x_l$  and  $f_q$  in terms of

remaining  $J$  variables which to be definite, we shall call  $F_r (r=1, 2, \dots, p)$  and  $X_s (s=p+1, \dots, J)$

Now we consider differentiable function  $g(x_j)$  as in [1]

$$g(x_j) = G(F_r, X_s) \quad (5)$$

at any point  $(F_r, X_s)$  in the vicinity of  $(F_r^0, X_s^0)$  such that  $F_r, f_q, x_l$  and  $X_s$  satisfy equation (2).

Let the value of  $G$  at the point  $(F_r^0, X_s^0)$  be designated by  $G^0$ . Using the approximation theorem about  $(F_r^0, X_s^0)$ , we have

$$\begin{aligned} g - G^0 &= \sum_{r=1}^p \left( \frac{\partial G}{\partial F_r} \right)^0 (F_r - F_r^0) + \sum_{s=p+1}^J \left( \frac{\partial G}{\partial X_s} \right)^0 (X_s - X_s^0) \\ &+ \frac{1}{2} \left[ \sum_{r=1}^p \left( \frac{\partial^2 G}{\partial F_r^2} \right)^0 (F_r - F_r^0)^2 + \sum_{\substack{l, k=1 \\ l \neq k}}^p \left( \frac{\partial^2 G}{\partial F_l \partial F_k} \right)^0 (F_l - F_l^0) (F_k - F_k^0) \right. \\ &+ \sum_{r=1}^p \sum_{s=p+1}^J \left( \frac{\partial^2 G}{\partial F_r \partial X_s} \right)^0 (F_r - F_r^0) (X_s - X_s^0) + \sum_{s=p+1}^J \left( \frac{\partial^2 G}{\partial X_s^2} \right)^0 (X_s - X_s^0)^2 \\ &\left. + \sum_{\substack{i, j=p+1 \\ i \neq j}}^J \left( \frac{\partial^2 G}{\partial X_i \partial X_j} \right)^0 (X_i - X_i^0) (X_j - X_j^0) \right] + O\{(F_r - F_r^0)^3, (X_s - X_s^0)^3\} \quad (6) \end{aligned}$$

The symbol  $\left( \frac{\partial G}{\partial F_k} \right)^0$  for some index  $k (1 \leq k \leq p)$  represents partial derivative of  $G$  with respect to  $F_k$  at  $(F_r^0, X_s^0)$ .

Now we have assumed that equations (2) are such that the transformed objective function  $G(F_r, X_s)$  is strictly convex differentiable function over the convex set "S" of feasible solutions  $x_j$ .

Therefore

$$\begin{aligned} L &= \frac{1}{2} \left[ \sum_{r=1}^p \left( \frac{\partial^2 G}{\partial F_r^2} \right)^0 (F_r - F_r^0)^2 + \sum_{r=1}^p \sum_{s=p+1}^J \left( \frac{\partial^2 G}{\partial F_r \partial X_s} \right)^0 (F_r - F_r^0) (X_s - X_s^0) \right. \\ &+ \sum_{\substack{l, k=1 \\ l \neq k}}^p \left( \frac{\partial^2 G}{\partial F_l \partial F_k} \right)^0 (F_l - F_l^0) (F_k - F_k^0) + \sum_{s=p+1}^J \left( \frac{\partial^2 G}{\partial X_s^2} \right)^0 (X_s - X_s^0)^2 \\ &\left. + \sum_{\substack{i, j=p+1 \\ i \neq j}}^J \left( \frac{\partial^2 G}{\partial X_i \partial X_j} \right)^0 (X_i - X_i^0) (X_j - X_j^0) \right] \quad (7) \end{aligned}$$

(@ If  $(f_i^0, x_j^0)$  is a basic feasible solution)

@ For definition of basic feasible, non-degenerate and relative minimum etc., reference may be had to Wilde's paper.

$$\begin{aligned}
 L = \frac{1}{2} \left[ \sum_{r=1}^p \left( \frac{\partial^2 G}{\partial F_r} \right)^o F_r^2 + \sum_{r=1}^p \sum_{s=p+1}^J \left( \frac{\partial^2 G}{\partial F_r \partial X_s} \right)^o F_r X_s + \sum_{\substack{l, k=1 \\ l \neq k}}^p \left( \frac{\partial^2 G}{\partial F_l \partial F_k} \right)^o F_l F_k \right. \\
 \left. + \sum_{s=p+1}^J \left( \frac{\partial^2 G}{\partial X_s^2} \right)^o X_s^2 + \sum_{\substack{i, j=p+1 \\ i \neq j}}^J \left( \frac{\partial^2 G}{\partial X_i \partial X_j} \right)^o X_i X_j \right] \quad (8)
 \end{aligned}$$

is a positive definite form, {[ 2 ] Appendix}

i. e.  $L > 0$ ,

for all values of  $F_r, X_s$  except when

$$\begin{aligned}
 F_r &= 0 \\
 X_s &= 0
 \end{aligned}$$

then  $L = 0$ . (9)

#### Theorem :

For a basic non-degenerate feasible solution to give a relative minimum, it is necessary and sufficient that

$$\left. \begin{aligned}
 \left( \frac{\partial G}{\partial F_r} \right)^o &\geq 0 & (r=1, \dots, p) \\
 \left( \frac{\partial G}{\partial X_s} \right)^o &\geq 0 & (s=p+1, \dots, J)
 \end{aligned} \right\} \quad (10)$$

The non-degeneracy of the solution implies that it will remain feasible for sufficiently small changes in  $F_r$  and  $X_s$ .

Necessary condition has been proved by Wilde in Lemma {1} of his paper.

To establish the sufficiency of these conditions we have to show that  $G_o$  is a relative-minimum, when

$$\begin{aligned}
 \left( \frac{\partial G}{\partial F_r} \right)^o &\geq 0 & (r=1, \dots, p) \\
 \left( \frac{\partial G}{\partial X_s} \right)^o &\geq 0 & (s=p+1, \dots, J)
 \end{aligned}$$

**Proof:** The proof is divided into two parts.

1st Part: At least one of the first order partial derivatives is not zero.

2nd Part: All the first order partial derivatives are zero.

#### FIRST PART

Let  $\left( \frac{\partial G}{\partial F_k} \right)^o > 0$  ( $1 \leq k \leq p$ ) while all other first order derivatives may be zero.

From (6)

$$g - G^0 = \left( \frac{\partial G}{\partial F_k} \right)^0 F_k + O(F_r^2, X_s^2) \quad (11)$$

A point  $(F_r, X_s)$  considered can be taken in such a small neighbourhood of  $(0, 0)$  i.e.  $(f_i^0, x_j^0)$  so that  $O(F_r^2, X_s^2)$  can be neglected in comparison to  $F_k \neq 0$

$$\begin{aligned} \text{therefore} \quad & g - G^0 > 0 \\ \text{or} \quad & g > G^0 \end{aligned}$$

In case  $F_k = 0$  then  $O(F_r^2, X_s^2)$  can no longer be neglected and then

$$g - G^0 = L + O(F_r^3, X_s^3) \quad (12)$$

The result is also obvious when more first order partial derivatives are not zero.

## SECOND PART

$$\left( \frac{\partial G}{\partial F_r} \right)^0 = 0 \quad (r = 1, \dots, p)$$

$$\left( \frac{\partial G}{\partial X_s} \right)^0 = 0 \quad (s = p+1, \dots, J)$$

Now

$$g(F_r, X_s) - G^0 = L + O(F_r^3, X_s^3) \quad (13)$$

Thus for (12) and (13) together, we have  $L$  as a positive definite form.

Therefore  $L > 0$  for all feasible  $(F_r, X_s)$  except when

$$\begin{aligned} F_r &= 0 & (r = 1, \dots, p) \\ X_s &= 0 & (s = p+1, \dots, J) \end{aligned}$$

then

$$L = 0$$

in that case our feasible solution coincides with  $(f_i^0, x_j^0)$  which is a *unique* basic feasible solution. Thus for all  $(F_r, X_s)$  other than  $(f_i^0, x_j^0)$

$$L > 0$$

$O(F_r^3, X_s^3)$  can be neglected in comparison to  $O(F_r^2, X_s^2)$  which occurs in  $L$ .

Therefore in some neighbourhood of  $(f_i^0, x_j^0)$

$$g - G^0 > 0$$

Thus  $G^0$  is a minimum (relative).

Hence for any feasible solution  $(f_i, x_j)$  in some neighbourhood of non-degenerate basic feasible solution  $(f_i^0, x_j^0)$

$$g > G^o$$

i. e.  $G^o$  is a relative minimum.

As in [1] we have

$$\begin{aligned} y_i &= \begin{cases} \frac{\partial G}{\partial f_i} & i=1, \dots, p \\ 0 & i=p+1, \dots, I \end{cases} \\ h_j &= \begin{cases} 0 & j=1, \dots, p \\ \frac{\partial G}{\partial x_j} & j=p+1, \dots, J \end{cases} \end{aligned} \quad (14)$$

Hence our theorem :

For a non-degenerate basic feasible solution  $(f_i^o, x_j^o)$  to give a relative-minimum (which is also absolute minimum) the necessary and sufficient conditions are

$$\begin{aligned} y_i^o &\geq 0 & i=1, \dots, I \\ h_j^o &\geq 0 & j=1, \dots, J \end{aligned} \quad (15)$$

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