

PLANAR BRANCHED COVERINGS OF GRAPHS

By

SHIGERU KITAKUBO

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Abstract: A graph \tilde{G} is called a *branched covering* of G with a *projection* $p: \tilde{G} \rightarrow G$ if there is a surjection $p: V(\tilde{G}) \rightarrow V(G)$ such that $p|_{N(\tilde{v})}: N(\tilde{v}) \rightarrow N(v)$ is surjective for any vertex $v \in V(G)$ and $\tilde{v} \in p^{-1}(v)$. It is said to be *regular* if there is a subgroup A of $\text{Aut}(\tilde{G})$ such that for any two vertices $u, v \in V(\tilde{G})$, $p(u) = p(v)$ if and only if $\tau(u) = v$ for some $\tau \in A$. In this paper we show that G has a planar and finite regular branched covering if and only if G is either planar or projective-planar.

1. Introduction

Let G be a simple, connected, finite graph and let $V(G)$ and $E(G)$ denote the vertex set and the edge set of G respectively. We call the set of vertices adjacent to a vertex $v \in V(G)$ the *neighborhood* of v and denote it by $N(v)$.

A graph \tilde{G} is called a *branched covering* of G with a *projection* $p: \tilde{G} \rightarrow G$ if there is a surjection $p: V(\tilde{G}) \rightarrow V(G)$ such that $p|_{N(\tilde{v})}: N(\tilde{v}) \rightarrow N(v)$ is a surjection for any vertex $v \in V(G)$ and $\tilde{v} \in p^{-1}(v)$. \tilde{G} is called a *covering* (or an *unbranched covering*) of G if $p|_{N(\tilde{v})}: N(\tilde{v}) \rightarrow N(v)$ is in particular a bijection. A branched covering is possibly infinite but is assumed to be connected throughout this paper.

A branched covering $p: \tilde{G} \rightarrow G$ is called *regular* if there is a subgroup A of $\text{Aut}(\tilde{G})$, the automorphism group of \tilde{G} , such that two vertices $u, v \in V(\tilde{G})$ project to the same vertex $p(u) (= p(v))$ of G if and only if some automorphism $\tau \in A$ carries u to v ($\tau(u) = v$). This group A is called the *covering translation group* of $p: \tilde{G} \rightarrow G$ and each element of A is called a *covering translation*. Note that the covering translations of a regular unbranched covering, except the identical translation, have no fixed point in $V(\tilde{G})$.

Our main result is the following.

Theorem 1.1. *A graph G has a regular branched covering \tilde{G} which is a finite planar graph if and only if G is either planar or projective-planar.*

This is motivated by a theorem recently proved by Negami [3].

Theorem 1.2 (Negami [3]). *A graph G has a regular covering \tilde{G} which is a finite planar graph if and only if G is either planar or projective-planar.*

If a graph G has a regular covering \tilde{G} , \tilde{G} is also a regular branched covering of G by definition. So Theorem 1.1 is an extension of Theorem 1.2. Combining this theorem with Corollary 2 in [2], we have the following.

Corollary 1.3. *If a graph G has a regular branched covering \tilde{G} which is a finite and planar graph, then G has a 2-fold planar unbranched covering.*

After some preliminaries in the next section, we prove the theorem, using Whitney's theorem [5] and the property of 2-orbifolds which can be found in Thurston [4, Chap. 13].

2. Lemmas

Denote the minimal degree of a graph \tilde{G} by $\delta(\tilde{G})$. Assuming $\delta(\tilde{G}) \geq 3$, we discuss the situation for a while. Let τ be any automorphism which belongs to a subset A of $\text{Aut}(\tilde{G})$. Suppose G is n -connected and U is an n -cut of G . A subset S of $V(G)$ is called a *fragment* for U if the induced subgraph $\langle S \rangle$ of G is a component of $G - U$. A fragment S is called an *A-equivariant fragment* if either $\tau(S) \cap S = \emptyset$ or $\tau(S) = S$.

The following lemmas are useful to prove the theorem when \tilde{G} is not 3-connected.

Lemma 2.1. *Let \tilde{G} be a connected, but not 3-connected, graph and $\delta(\tilde{G}) \geq 3$. If $p: \tilde{G} \rightarrow G$ is a finite regular branched covering with covering translation group A of order ≥ 2 , then the following three hold:*

- (1) *The inclusion minimal fragment S for a 1- or 2-cut U of \tilde{G} is an A -equivariant fragment.*
- (2) *If $K = \langle S \cup U \rangle$ is not 3-connected, then U is a 2-cut, say $\{u, v\}$, and $K + uv$ is 3-connected.*
- (3) *If $p(U)$ is not a cutset of G , $V(G) = p(S) \cup p(U)$.*

Proof. Choose a 1- or 2-cut U of \tilde{G} and the fragment S for U such that S is minimal with respect to inclusion among all the fragments of \tilde{G} . Note that S contains neither a 1-cut nor a vertex of any 2-cut of \tilde{G} .

Let $\tau \in A$ be any covering translation. Since τ is an automorphism of \tilde{G} , $\tau(U)$ is also a cutset of \tilde{G} and $|U| = |\tau(U)|$. So we consider how U and $\tau(U)$ are placed in \tilde{G} . Considering the minimality of S and the property of U , as mentioned above, we can decide the location of the vertices and we get seven cases shown in Fig. 2.1. For example, in the first case, that is, when U is a

1-cut and $\tau(U) \neq U$, $\tau(S) \cap S = \emptyset$ because \tilde{G} is connected.

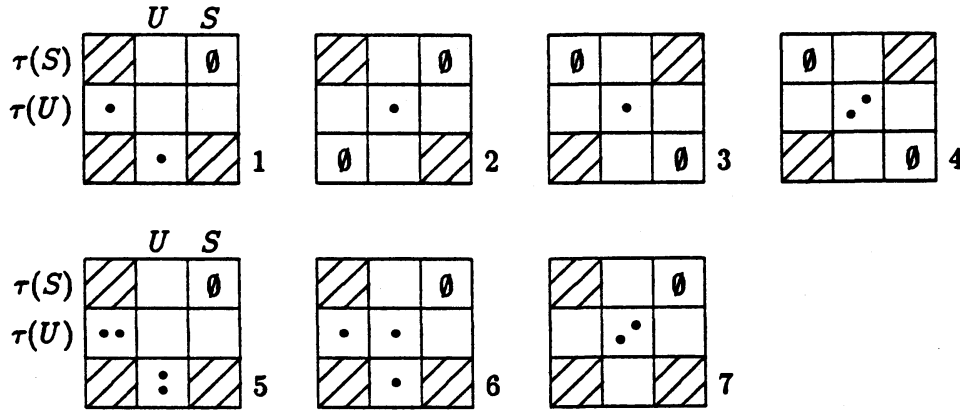


Fig. 2.1.

In each of these cases either $\tau(S) \cap S = \emptyset$ or $\tau(S) = S$. Thus (1) follows. Note that neither of the vertices in $\tau(U)$ lies in S .

Let $K = \langle S \cup U \rangle$ with S being the same minimal fragment as chosen above and suppose K has either a 1- or 2-cut U_K .

If $U = \{u\}$, U_K is also a cutset of \tilde{G} . This contradicts the minimality of S . So U must be a 2-cut $\{u, v\}$. Let $U_S = S \cap U_K$. Since $\langle S \rangle$ is connected, $U_S \neq \emptyset$ and U_S is not a cutset of \tilde{G} by the minimality of S . If U_S is a 1-cut u' , $\{u, u'\}$ is a 2-cut of \tilde{G} . This contradicts the minimality of S , so $U_S = U_K$.

Let the components of $\langle S - U_S \rangle$ be W_1, W_2, \dots, W_k . Then for $i = 1, 2, \dots, k$, W_i has at least one vertex which belongs to either $N(u)$ or $N(v)$ in K because of the minimality of S . So U_S is not a cutset of $K + uv$. Thus (2) follows.

Choose a vertex $\tilde{y} \in S$ and fix it. Set $y = p(\tilde{y})$. Assume that $G - p(U)$ is connected. Then there is a $y-x$ path Q_x in $G - p(U)$ for any vertex $x \in V(G) - p(U)$. Starting from \tilde{y} and naturally tracing edges in \tilde{G} which project to the edges on Q_x , we can get the path lift \tilde{Q}_x of Q_x and its terminal vertex \tilde{x} projects to x . Since Q_x does not meet $p(u)$ and $p(v)$, its lift \tilde{Q}_x is a path in $\langle S \rangle$. Therefore, for any vertex $x \in V(G) - p(U)$, there is a vertex $\tilde{x} \in S$ such that $p(\tilde{x}) = x$. This implies that $V(G) = p(S) \cup p(U)$. Thus (3) follows. \square

Lemma 2.2. Let \tilde{G} , p , and U be as in Lemma 2.1. If \tilde{G} is a planar branched covering with $|V(\tilde{G})|$ minimal, $p(U)$ must be a cutset of G .

Proof. Suppose $p(U)$ is not a cutset of G by Lemma 2.1, \tilde{G} has an A -equivariant fragment S and $V(G) = p(S) \cup p(U)$.

Here we consider a subgroup B of A which acts as an automorphism group of $K = \langle S \cup U \rangle$, that is, $A \supset B = \{\tau \mid \tau(K) = K\}$.

1) When $U = \{u\}$, let $K' = K$.

2) When $U = \{u, v\}$ and $\tau(u) = v$ for some $\tau \in B$, let $K' = K$.

3) When $U=\{u, v\}$ and $\tau(u)=v$ for some $\tau \in A-B$, let K' be the graph obtained from K by identifying u with v .

4) When $U=\{u, v\}$ and $\tau(u) \neq v$ for all $\tau \in A$ and $uv \notin E(\tilde{G})$ and $p(u)p(v) \in E(G)$, let $K'=K+uv$.

5) When $U=\{u, v\}$ and $\tau(u) \neq v$ for all $\tau \in A$ and $uv \notin E(\tilde{G})$ and $p(u)p(v) \notin E(G)$, let $K'=K$.

6) When $U=\{u, v\}$ and $\tau(u) \neq v$ for all $\tau \in A$ and $uv \in E(\tilde{G})$ and $p(u)p(v) \in E(G)$, let $K'=K$.

In each case of 1)~6), it is clear that we still have the same quotient $K'/B=G$. So K' is a regular branched covering of G , which contradicts the minimality of $|V(\tilde{G})|$. \square

3. Proof of the theorem

Suppose that \tilde{G} is a finite regular branched covering of G and that \tilde{G} is planar. Let $p: \tilde{G} \rightarrow G$ be the branched covering projection and $A \subset \text{Aut}(\tilde{G})$ the covering translation group.

When \tilde{G} is 3-connected, the theorem follows fairly immediately from Whitney's theorem [5]:

Theorem 3.1 (Whitney [5]). *A 3-connected planar graph can be embedded in a sphere such that each automorphism of the graph extends to a homeomorphism of the sphere onto itself.*

We call this embedding a *faithful embedding* [1].

Embed \tilde{G} faithfully in a sphere S^2 , and the translation group $A \subset \text{Aut}(\tilde{G})$ can be regarded as a group consisting of some self-homeomorphism $h: S^2 \rightarrow S^2$ such that $h(\tilde{G})=\tilde{G}$. Thus we can embed \tilde{G}/A in S^2/A .

So the following diagram holds.

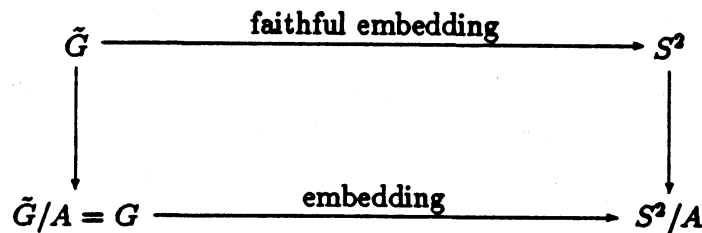


Fig. 3.1.

S^2/A is the quotient of S^2 by the action of A , called an elliptic 2-orbifold in [4], and its underlying space is homeomorphic to S^2 , D^2 or P^2 . Since $G=\tilde{G}/A$, G can be embedded in a sphere or a projective plane.

This completes the proof when \tilde{G} is 3-connected.

The rest of this paper is for the proof when \tilde{G} is not 3-connected. We use induction on the number of vertices of G .

When $|V(G)| \leq 6$, G is either planar or projective-planar. For the complete graph K_6 can be embedded in a projective plane, as depicted in Fig. 3.2.

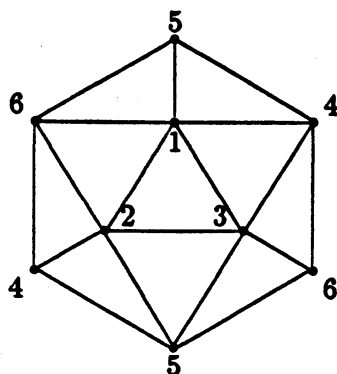


Fig. 3.2.

Next we divide the case into two; when $\delta(\tilde{G}) \geq 3$ and when otherwise. When $\delta(\tilde{G}) \geq 3$, we suppose any graph with finite planar branched covering that has fewer vertices than G is either planar or projective-planar.

First we modify $K = \langle S \cup U \rangle$ and $p(K)$ to prove that $p(K)$ is planar.

(1) When $U = \{u\}$, let $K' = K$ and $G_2 = p(K)$. Then $p' = p|_{K'} : K' \rightarrow G_2$ is a branched covering with the covering translation group $B = \{\tau|_{K'} : \tau(K) = K, \tau \in A\}$. By Lemma 2.1, K' is 3-connected and so K' can be faithfully embedded in the sphere S^2 by Theorem 3.1. Since $\tau(U) \cap S = \emptyset$ for each $\tau \in B$ (this is obvious when we check the shape of \tilde{G} described in the proof of Lemma 2.1), u is a fixed point of all $\tau \in B$ and then S^2/B is an elliptic 2-orbifold with elliptic points. The underlying space of such 2-orbifolds is homeomorphic to D^2 or S^2 and so G_2 is planar.

(2) When $U = \{u, v\}$ and $\tau(u) \neq v$ for all $\tau \in A$, let $K' = K + uv$ and $G_2 = p(K) + p(u)p(v)$. Then $p' = p|_{K'} : K' \rightarrow G_2$ is a branched covering with the covering translation group $B = \{\tau|_{K'} : \tau(K) = K, \tau \in A\}$. By Lemma 2.1, K' is 3-connected and so K' can be faithfully embedded in the sphere S^2 by Theorem 3.1. Since $\tau(u) \notin S$ for each $\tau \in B$, uv is a set of fixed points with respect to B and then S^2/B is an elliptic 2-orbifold with a mirror. The underlying space of such 2-orbifolds is homeomorphic to D^2 and so G_2 is planar.

(3) When $U = \{u, v\}$ and $\tau(u) = v$ for some $\tau \in A$ and $\tau(u) \neq v$ for all $\tau \in B$, let $K' = K + uv$ and G_2 be the graph obtained from $p(K)$ by splitting the vertex $p(u) (= p(v))$. (Fig. 3.3 shows an example of splitting a vertex.) Then $p' = p|_{K'} : K' \rightarrow G_2$ is a branched covering with the covering translation group $B = \{\tau|_{K'} : \tau(K) = K, \tau \in A\}$. By Lemma 2.1, K' is 3-connected and so K' can be faithfully embedded in the sphere S^2 by Theorem 3.1. Since $\tau(u) \notin S$ for each

$\tau \in B$, uv is a set of fixed points with respect to B and then S^2/B is an elliptic 2-orbifold with a mirror. The underlying space of such 2-orbifolds is homeomorphic to D^2 and so G_2 is planar. Since G_2 was obtained from $p(K)$ by splitting the vertex $p(u)$, $p(K)$ is clearly planar.

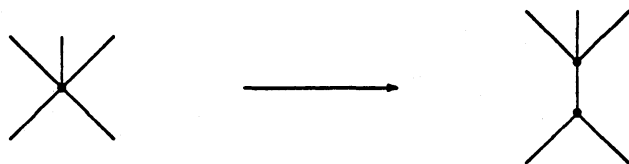


Fig. 3.3.

(4) When $U = \{u, v\}$ and $\tau(u) = v$ for some $\tau \in B$, let $G_2 = p(K)$. Then $p' = p|_K: K \rightarrow G_2$ is a branched covering with the covering translation group $B = \{\tau|_K: \tau(K) = K, \tau \in A\}$. By Lemma 2.1, $K + uv$ is 3-connected and B is still an automorphism group of $K + uv$, so we embed $K + uv$ faithfully in the sphere S^2 . Since $\tau(u) \notin S$ for each $\tau \in B$, the middle point of uv is the only fixed point on uv . In the same manner as the case (1), $(K + uv)/B$ is planar. Then clearly G_2 is planar.

In any case G_2 is planar. Let $G_1 = G - p(S) + p(u)p(v)$. If $p(u) = p(v)$, let $G_1 = G - p(S)$. Then $G = G_1 \cup G_2$ or $G_1 \cup G_2 - p(u)p(v)$.

By induction hypothesis, G_1 is either planar or projective-planar. Therefore we can easily embed G in a plane or projective plane.

This completes the proof when $\delta(\tilde{G}) \geq 3$.

When \tilde{G} is not 3-connected and $\delta(\tilde{G}) < 3$, we will show that if a planar graph \tilde{G} is a regular branched covering of G , the graph \tilde{G}' , which is obtained from \tilde{G} by deleting all the vertices that have degree 1 or 2 in \tilde{G} , is also a regular branched covering of G' which is obtained from G by deleting all the images of the deleted vertices of \tilde{G} . Then we can reduce this case to the previous case.

Let $p' = p|_{\tilde{G}}$, and $A' = \{\tau|_{\tilde{G}}: \tau \in A\}$ in the following arguments (1) and (2), and we show how to get \tilde{G}' and G' .

(1) Suppose that $\deg_{\tilde{G}}(u) = 1$ for some vertex $u \in V(\tilde{G})$.

Let \tilde{G}' be

$$\tilde{G} - \bigcup_{\tau \in A} \tau(u)$$

and G' be $G - p(u)$. Then \tilde{G}' is planar and is a regular branched covering of $G' = \tilde{G}'/A$ as depicted.

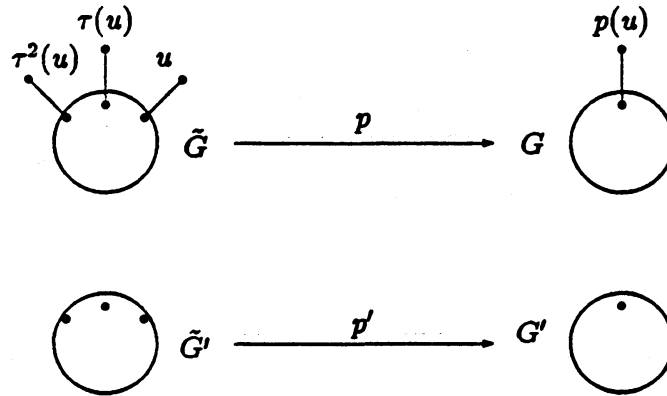


Fig. 3.4.

(2) Suppose that $\deg_{\tilde{G}}(u)=2$ and u is adjacent to x and y .

a) If x is not adjacent to y and if $\tau(x) \neq y$ for any $\tau \in A$, let \tilde{G}' be

$$\tilde{G} - \bigcup_{\tau \in A} \tau(u) + \bigcup_{\tau \in A} \tau(x)\tau(y)$$

and G' be $G - p(u) + p(x)p(y)$.

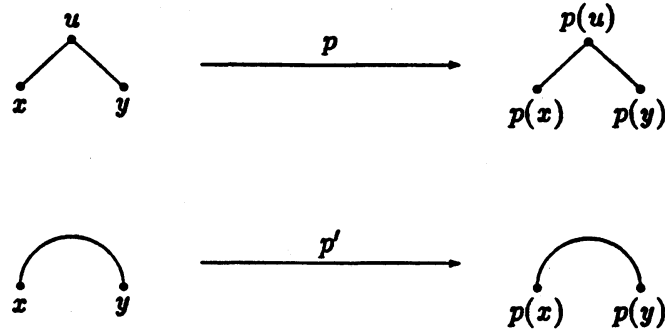


Fig. 3.5.

b) If x is not adjacent to y and if $\tau(x)=y$ for some $\tau \in A$, let \tilde{G}' and G' be the same as (1).

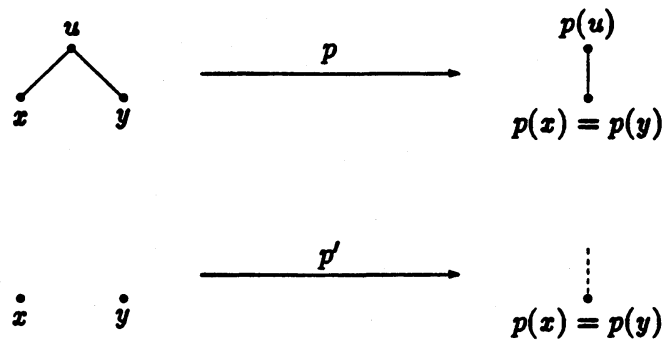


Fig. 3.6.

c) If x is adjacent to y , let \tilde{G}' and G' be also the same as (1).

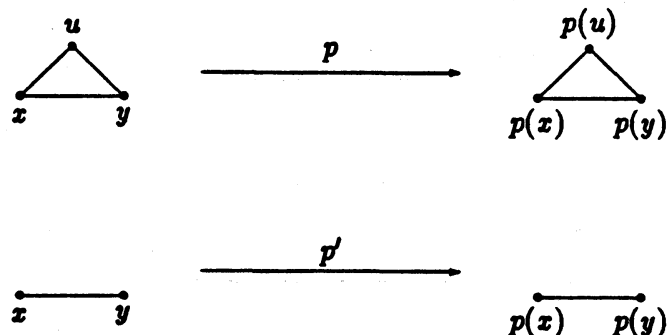


Fig. 3.7.

In each case of a), b), and c), \hat{G}' is a regular branched covering of G' .

By induction hypothesis G' is either planar or projective-planar. Clearly we can get a planar or projective-planar embedding of G by adding the vertex $p(u)$, which has degree at most 2.

This completes the proof when $\delta(\tilde{G}) < 3$.

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Department of Information Science
Tokyo Institute of Technology
Oh-Okayama, Meguro-ku,
Tokyo, 152, Japan